

Towards Automatically Estimating Porting Effort Between Web Service APIs

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ABSTRACT

Developers today increasingly incorporate curated web services, accessed over a network via well-defined and published interfaces (APIs), as modules in their applications. Public versions of these *web APIs* emerge and change frequently, making it critical for software development, testing, and maintenance personnel to be able to estimate the workload associated with “porting” (or migrating) an application to a new API or API version. Unfortunately, today there is no simple automated mechanism for estimating and reasoning about the application porting effort that will be necessary when the web APIs that an application uses change.

To address this limitation, we describe an automated methodology for quantifying the porting effort associated with the use of web APIs. Our approach defines a simple language (based on Python) with which API developers specify the semantics of API operations, a tool set that consumes and extracts semantic similarity of API operations from annotations expressed in this language, and a metric that facilitates ranking of porting effort for API operation pairs. We evaluate our approach using both randomly generated and real-world APIs and show that our metric can correctly categorize the relative difficulty that developers associate with porting an application from one API to another.

Categories and Subject Descriptors

D.2.2 [Software Engineering]: Design Tools and Techniques—*Modules and interfaces*; D.2.7 [Software Engineering]: Distribution, Maintenance, and Enhancement—*Portability*

General Terms

Experimentation, Languages, Management

Keywords

Web services, Web APIs, Porting effort, Semantic similarity, Axiomatic semantics

1. INTRODUCTION

Web services are used increasingly for the implementation of Internet accessible applications [1, 2, 3]. In this emerging development model, application architects and programmers use extant, network accessible services as external modules in their applications (or, indeed, as parts of other web services). Building applications from curated web services improves programmer productivity over non-service-oriented methodologies by easing assembly, testing, maintenance, and by improving the robustness of complex systems through the reuse of software and data components offered by providers “as-a-service” [4]. By composing an application from existing services that encapsulate common yet complicated tasks (data access and analysis, messaging, logging, security, etc.), application developers are able to work at a higher level of abstraction, thereby saving valuable development and debugging time, while at the same time providing scalable access to the core of the application logic via a network-facing service deployment. In addition, the services that an application calls upon are typically maintained and curated by an IT staff which is responsible for assuring service integrity. That is, an application that composes running web services, leverages the deployment and maintenance resources devoted to those services.

A web service itself consists of one or more software components each with a well-defined but separate (in terms of coding and implementation) application programming interface (API) that is network-accessible and facilitates machine-to-machine interoperation. The web service “stack” is responsible for connecting each service implementation to the API code that exposes it to its users.

The growth in the popularity of this approach to application development has introduced several challenges for developers. In particular, because web-service-based applications decouple their service implementations from their APIs, the development and maintenance life cycles for APIs and service implementations are separated. As a result, APIs can change independently of the implementations they serve and new APIs (offering additional features as a superset) emerge frequently for existing services.

For example, commercial service providers experience com-

petitive pressure to add, modify, deprecate, and retire APIs. New APIs are frequently introduced by competitors that are similar in functionality to existing APIs but offer some additional functional and/or business advantage. As a result, many commercial service providers change their APIs in terms of their operations, functionality, and behavior as well as the terms of their licensing, service level agreements (SLAs), and pricing, over time.

There are several notable examples of such commercial API churn. Amazon produced 86 releases of the Amazon EC2 APIs between August 2006 to August 2013 [5]. Eucalyptus, which mirrors the Amazon EC2 API in open source for free, on-premise use of Amazon-compatible services [6], releases its software quarterly to keep pace. Twitter released the version 1.1 of their web API on September 2012 and pulled version 1.0 out of production on May 2013 [7]. eBay has released 13 versions of its trading API during the first two quarters of year 2013 [8]. The licensing terms of the Amazon Product Advertising API changed twice between 2011 and 2013 [9].

In this paper, we describe a methodology for quantifying the effort required to port an application from one API to another, or from one API version to another version of the same API. With the possibility of high churn, it is critical that the software development, testing, and maintenance activities be informed by measures of the difficulty associated with “porting” an application to a new API or API version. We refer to this measure, in the abstract, as *porting effort*. Porting effort includes the effort needed to rewrite and refactor code, perform regression testing, and update all impacted test cases, configurations, deployment scripts, and documentation. At present, we know of no simple mechanism for estimating and reasoning about the amount of “work” required to migrate an application between web service APIs. The state of the practice is that, developers simply perform the migration speculatively (possibly losing their work if the migration fails) or rely on their intuition and experience to obtain this measure.

To address this limitation, we propose a simple but formal mechanism to estimate porting effort for applications that use web APIs. Our approach focuses on semantic similarity of APIs and trades off complexity with efficacy and speed of analysis. Our goal is to capture the functional behavior of API operations (using existing, well-understood techniques) with sufficient accuracy and minimal programmer intervention to rank multiple API alternatives by porting effort.

Note that, in this work, we do not address the research question associated with ensuring that an API and its service implementation remain consistent and conforming to the specification. That is, because the API and implementation are essentially developed and managed as separate software components with separate life cycles, it is possible for their respective functionalities to fail to meet the specification when they are composed. The current state of the practice puts the burden for ensuring API and service implementation conformity on the development and/or maintenance staff. Indeed, the concept of DevOps [10] has been developed, in part, to address this problem organizationally. We believe that future work building on the large

body of results from program verification research will address this worthy topic. In this work, however, we look only at the question of compatibility and consistency between APIs themselves and not their respective service implementations.

Initially we focus on a straightforward application of program semantics applied to the APIs of web services. This starting point exposes the working of our mechanism while providing a framework that can make use of more advanced techniques from the program verification research corpus [11, 12, 13, 14, 15] and semantic web technologies [16, 17, 18, 19, 20, 21, 22, 23].

To enable this investigation, we define a simple language in terms of a restricted subset of the popular Python programming language¹. Developers use this language to specify the axiomatic semantics [24] of web APIs (*i.e.* the preconditions and postconditions associated with each operation). Our Python-based specification language is familiar to many developers while facilitating simple static analysis. We use the output of this analysis to measure API similarity via an extended form of the Dice coefficient [25] on the abstract syntax trees of semantic predicates, combined with Hoare’s consequence rule [24] applied to API pairs.

We implement our approach and evaluate it using a number of popular APIs for social media login, airline itinerary search, and digital media video search. Our initial results indicate that developers can quantify and reason about the porting effort of migrating their applications to different web API versions and competitive implementations, without performing the porting. Our experimental results also show our approach to be efficient enough to be a practical part of the software engineering process used to develop service-composing applications. In the sections that follow, we overview porting effort and detail our approach. We then describe the empirical evaluation of our approach, discuss the results, and conclude.

2. EVALUATING PORTING EFFORT

We start with the hypothesis that application porting effort from one API to another is correlated with the degree to which two APIs are similar. Two APIs are comparable in terms of porting effort if they are two different versions of the same API or expose same or similar services. API similarity can be syntactic, *i.e.*, two APIs export operations with similar cardinality and data types for their inputs and outputs. Alternatively, similarity can be semantic, *i.e.*, two APIs are similar in terms of the functionality and behavior of their syntactically similar operations. In this work, we focus on the latter, *i.e.* on the functional similarity of the operations of two APIs that match syntactically. We plan to consider syntactic similarity more comprehensively in future work.

To define a metric for application porting effort from one API to another using the semantics of their operations, we require mechanisms

- with which API developers specify the semantics of

¹Python is used heavily in web services, particularly as a prototyping language.

API operations (Subsection 2.1) ²,

- that automate the consumption and analysis of specified API semantics (Subsection 2.2), and
- that use the output from the analysis to construct a measure of porting effort for a pair of APIs (Subsection 2.3).

To define each of these mechanisms and the overarching metric, we leverage and assemble extant research advances in a simple, yet new way that enables developers to estimate and rank the effort associated with porting their application to a different version of a web API or to an alternative implementation of an API. For simplicity of discussion, we assume that a pair of APIs under consideration have a *single*, syntactically matching operation. That is, in what follows, we will examine the ability to quantify similarity between individual API operations. As part of our future work we plan to extend the methodology to consider multiple operations in pairs of APIs.

2.1 Specifying Web API Semantics

The first mechanism of our approach is a specification language that developers can use to document the semantics of the operations in their web APIs. Our goal is to define a language that is simple, familiar, and intuitive to use that, at the same time, enables developers to specify the meaning of an API in a way that is amenable to efficient static analysis for semantic similarity. Toward this end, we leverage popular programming language syntax and tooling, and the well researched field of axiomatic semantics.

Our language is a strict subset of the Python programming language. This language choice is inspired by the widespread use of Python, Python’s high level of abstraction and available tooling, and by previous works such as JML [26] and Spec# [27] that document program semantics (behavioral interface specifications) using programming language syntax. This latter research and that of others shows that using the syntax of familiar and popular programming languages to document API semantics facilitates programmer creation and editing of semantic specifications [28].

We restrict the Python language in a number of ways to facilitate analysis and to simplify the specification process by API developers. Our language only accepts single-lined Python statements that are free of side effects. We disallow side effects to preclude the consideration of internal service state. We also disallow conditionals, loops, try-catch blocks, class definitions, and function definitions.

Developers use this language to describe the behavior of API operations using axiomatic semantics – preconditions that hold prior to invoking the operation and postconditions that hold after the operation executes. We leverage axiomatic semantics as a first step toward describing and analyzing API operations in a way that reflects porting effort. We plan to consider other successful approaches [29, 30, 11] to describing the function and behavior of API operations as part of future work.

²We assume that the service provider has verified that the API semantics match those of the service implementation as part of its quality assurance and testing process.

Developers refer API request parameters and response parameters using the built-in logical variables `input` and `output`, respectively. For example, for an operation that takes two positive numbers and responds with their sum, the preconditions can be documented using the statements `input.x > 0` and `input.y > 0`; the postconditions can be documented as `output.sum == input.x + input.y`. These logical variables have been inspired by Hoare logic [24] and separation logic [31] to differentiate precondition values from postcondition values. The use of logical variables also enables expressing postconditions relative to preconditions, that is, postconditions can refer to the pre-state (request state) of an operation.

We do not allow invoking arbitrary functions using our language. This includes the built-in functions of Python as well as any class-level functions that can be invoked as object methods. However, we do support a number of useful predefined, side-effect-free, functions (that we have defined) when invoked as built-in functions (as opposed to object methods). We currently support the functions `len`, `implies`, `forall`, `exists`, `matches`, `datebefore`, and `dateformat`. We illustrate the use of a subset of our built-ins using simple examples below. Our language and built-ins are easily extended if and when more expressive power is required.

- The `password` input must be at least 6 characters long:
`len(input.password) >= 6`
- All entries in the input list named `scores` must be within the range [0,100]:
`forall(entry, input.scores, 0 <= entry and entry <= 100)`
- The format of the `publishedDate` output field is `yyyy-MM-dd`:
`dateformat(output.publishedDate, 'yyyy-MM-dd')`
- If the `country` input field is set to `US`, the `currency` output field will be set to `USD`:
`implies(input.country == 'US', output.currency = 'USD')`

2.2 Comparing API Operations Pairwise

We next determine a similarity “score” by comparing the preconditions and postconditions of individual API operations. Throughout the remainder of this paper, we refer to the specified preconditions and postconditions of an API simply as *semantic predicates*. We represent semantic predicates as abstract syntax trees (ASTs).

To compare a pair of matching API operations, we compute a tree similarity metric on their ASTs. To enable this, we employ a technique that is widely used for software plagiarism detection and source code evolution analysis, called the Dice coefficient [32]. The Dice coefficient has been shown in this past work to accurately extract the semantic similarity of two code fragments. Using the Dice coefficient, we treat each AST as a set of nodes over which we compute set similarity. Specifically, if P_1 and P_2 are two semantic predicates whose ASTs are T_1 and T_2 respectively, we compute the degree of similarity between the predicates P_1 and P_2 by computing the Dice coefficient on T_1 and T_2 as follows.

$$\text{Similarity}(\langle P_1, P_2 \rangle) = \text{Dice}(\langle T_1, T_2 \rangle) \quad (1)$$

$$Dice(\langle T_1, T_2 \rangle) = \frac{2C}{2C + L + R} \quad (2)$$

C is the number of nodes common to both T_1 and T_2 . L is the number of nodes unique to T_1 and R is the number of nodes unique to T_2 . This approach enables us to obtain a similarity value between 0 and 1 for any two given semantic predicates, where 0 indicates a total mismatch and 1 indicates a perfect match.

We also apply a trivial transformation on the semantic predicates when performing semantic comparison that breaks disjunctive and conjunctive predicates into their constituent predicates. This enables our mechanism to handle situations where the same set of predicates have been expressed in two APIs, but in slightly different formats.

Notice that the amount of work necessary to port from one API to another is affected by the number of predicates in each. In particular, the effort to port from a source API with fewer preconditions than the target API is more difficult than porting in the reverse direction.

To illustrate this asymmetry, let M and N be two web APIs where N has more preconditions than M . It is more difficult to port from M to N than from N to M . More preconditions imply that N 's input set is more restricted than M . Therefore it cannot support all the inputs that M does. Hence some extra effort has to be put in by the developer to make sure that the application doesn't pass an unsupported input value to API N . However, by the same argument, porting an application from N to M should be easier. Since M 's input set is less restricted than N , the developer doesn't have to do any extra work in this case.

Notice also that a similar asymmetry exists with respect to postconditions. If an application is to be ported from API S to API T and if T has more postconditions than S , then porting S to T is easier than the other way around. More postconditions help further restrict the output of API T . In other words, T may not produce an output that S doesn't. Therefore the application should be able to handle all the outputs generated by T , without having to make any code changes. On the other hand, porting from API T to S becomes more difficult, since S might produce an output that T doesn't.

2.3 Quantifying Porting Effort Between API Operations

Using the mechanism described in the previous section, we construct a measure of application porting effort using the semantic similarity of two APIs. Suppose S is a source API with the precondition set S_{pre} and the postcondition set S_{post} . Suppose T is a target API with the precondition set T_{pre} and the postcondition set T_{post} . To compute the porting effort from S to T , we first compare each member in S_{pre} against each member in T_{pre} . That is, we calculate the similarity (Dice coefficient) of each predicate pair in $S_{pre} \times T_{pre}$. Then we choose the pairs with the highest similarity, and match each member in S_{pre} to a member in T_{pre} . In other words, for each predicate $x \in S_{pre}$ we assign a predicate $y \in T_{pre}$ such that the similarity of $\langle x, y \rangle$ is greater than

the similarity of any $\langle x, z \rangle$ where $z \in T_{pre}$ and $y \neq z$. Matched pairs are put into a new set M_{pre} . We also make sure that no member in S_{pre} or T_{pre} is matched to multiple counterparts. That is, whenever we insert a pair $\langle x, y \rangle$ into M_{pre} , we mark x in S_{pre} and y in T_{pre} so that they cannot be considered for a match again. This way each member in S_{pre} can be matched to a unique member in T_{pre} as long as $|S_{pre}| \leq |T_{pre}|$. But if $|S_{pre}| > |T_{pre}|$ some members of S_{pre} will remain unmatched.

We translate the predicate assignments into a porting effort score by computing $(1 - D_i)$ where D_i is the similarity of the pair $i \in M_{pre}$. We add these values up to obtain an initial porting effort score P_{eff1} . Then we consider the remaining unmatched (unmarked) predicates in S_{pre} and T_{pre} . Recall that porting to an API with more preconditions is more difficult than in the reverse direction. To reflect this asymmetry in our methodology, we increase P_{eff1} by 1 for each unmatched predicate in T_{pre} . Unmatched predicates in S_{pre} are ignored. Therefore, we have:

$$P_{eff1}(S, T) = \sum_{i \in M_{pre}} (1 - D_i) + |T_{pre}| - |M_{pre}| \quad (3)$$

We perform a similar computation for postconditions using the sets S_{post} and T_{post} . We compute the similarity of the members of $S_{post} \times T_{post}$ and pick the pairs with the highest similarity to initialize a matching set M_{post} . As a postcondition pair $\langle x, y \rangle$ inserted to M_{post} , we mark x in S_{post} and y in T_{post} to ensure that no predicate is matched multiple times. Then for each pair $j \in M_{post}$ we compute $(1 - D_j)$ where D_j is the similarity of the pair j , and add these values up to obtain the porting effort score P_{eff2} . We further penalize the porting effort by increasing P_{eff2} by 1 for each unmarked (unmatched) predicate in S_{post} . This adjustment accounts for the greater difficulty associated with porting from an API with more postconditions to one with fewer postconditions.

$$P_{eff2}(S, T) = \sum_{j \in M_{post}} (1 - D_j) + |S_{post}| - |M_{post}| \quad (4)$$

We calculate the final porting effort score by combining the values obtained from previous computations. If $P_{eff}(S, T)$ is the overall porting effort from API S to API T , we have:

$$P_{eff}(S, T) = P_{eff1}(S, T) + P_{eff2}(S, T) \quad (5)$$

Algorithm 1 further illustrates our porting effort evaluation method. $Temp_1$ and $Temp_2$ are map data structures that support storing key-value pairs. The algorithm makes use of following named procedures:

- *map_store*(*map*, *key*, *value*) - Stores the given *key* - *value* pair in the *map*.
- *map_get_max*(*map*) - Returns the key-value pair with the largest value in the *map*.
- *map_remove*(*map*, *key*) - Removes the entry with the specified *key* from the *map*.

- $mark(set, element)$ - Marks the specified *element* in the *set*.
- $unmarked(set)$ - Returns *TRUE* if the *set* contains at least one unmarked element. Otherwise returns *FALSE*.
- $Sim(< x, y >)$ - Returns the similarity (Dice coefficient) of the predicate pair $< x, y >$.

```

Data: Source API  $S$  with predicate sets  $S_{pre}, S_{post}$  and
        Target API  $T$  with predicate sets  $T_{pre}, T_{post}$ 
Result: Porting effort
 $M_{pre} \leftarrow \emptyset, M_{post} \leftarrow \emptyset$ 
 $P_{eff1} \leftarrow 0, P_{eff2} \leftarrow 0$ 
 $Temp_1 \leftarrow EmptyMap, Temp_2 \leftarrow EmptyMap$ 
for  $< x, y > \in (S_{pre} \times T_{pre})$  do
  |  $map\_store(Temp_1, < x, y >, Sim(< x, y >))$ 
end
while  $unmarked(S_{pre})$  and  $unmarked(T_{pre})$  do
  |  $<< x, y >, D_i > \leftarrow map\_get\_max(Temp_1)$ 
  |  $mark(S_{pre}, x), mark(T_{pre}, y)$ 
  |  $map\_remove(Temp_1, < x, y >)$ 
  |  $M_{pre} \leftarrow M_{pre} \cup \{< x, y >\}$ 
  |  $P_{eff1} \leftarrow P_{eff1} + (1 - D_i)$ 
end
 $P_{eff1} = P_{eff1} + |T_{pre}| - |M_{pre}|$ 
for  $< x, y > \in (S_{post} \times T_{post})$  do
  |  $map\_store(Temp_2, < x, y >, Sim(< x, y >))$ 
end
while  $unmarked(S_{post})$  and  $unmarked(T_{post})$  do
  |  $<< x, y >, D_j > \leftarrow map\_get\_max(Temp_2)$ 
  |  $mark(S_{post}, x), mark(T_{post}, y)$ 
  |  $map\_remove(Temp_2, < x, y >)$ 
  |  $M_{post} \leftarrow M_{post} \cup \{< x, y >\}$ 
  |  $P_{eff2} \leftarrow P_{eff2} + (1 - D_j)$ 
end
 $P_{eff2} = P_{eff2} + |S_{post}| - |M_{post}|$ 
return  $P_{eff1} + P_{eff2}$ 

```

Algorithm 1: Porting effort evaluation algorithm

3. EXPERIMENTAL RESULTS

We have implemented a prototype of our methodology that takes two web API specifications as the input, and outputs the porting effort between them. The input web API specifications are simple JSON strings, similar to Swagger API descriptions [33]. The specifications list all of the operations of the web APIs along with their axiomatic semantics expressed using Python-based specification language described in subsection 2.1.

We use this prototype to investigate the properties of our methodology as well as to expose its efficacy. In the first set of experiments, we consider randomly generated API specifications to study various characteristics of our API porting effort metric. We then consider real-world APIs and developer-perceived porting effort, and evaluate the overhead of API similarity mechanism.

3.1 Randomly Generated APIs

In our first experiment, we randomly generate a population of 100 API specifications. Each specification has a single op-

eration. We compare each API against all others in the population and compute the porting effort between them. We repeat this experiment using different numbers of semantic predicates. We randomly generate the API specifications with 10, 20 and 50 semantic predicates. Our goal with this experiment is to understand how our measure of porting effort changes under these scenarios (*e.g.* to determine the sensitivity of the mechanism to supplied parameters).

Figure 1 shows the cumulative distribution functions (CDFs) of the computed porting effort as a function of the number of predicates per single API operation. A porting effort value of 0 indicates no porting effort. The data shows that the porting effort between API operations increases with the number of semantic predicates. For example, the maximum porting effort observed in APIs with 10 semantic predicates is 17.4. This goes up to 30.1 when the number of predicates is increased to 20. It further increases up to 44.9 when the semantic predicates count is set to 50. Also, when considering the CDFs of the porting effort, 50% of the API operation pairs have 4.3 or less porting effort in the population with 10 semantic predicates. In the population with 20 semantic predicates, 50% of the APIs have 7.1 or less porting effort. In the population with 50 semantic predicates, this limit further increases up to 12.9. This is inline with our experience in which, as the number of semantic predicates increases, the API consumer is forced to adhere to additional restrictions. As such, when porting among different web API operations, the developer has to take more constraints into account and must write more code to reconcile the differences. This results in increased porting effort. Our experimental results suggest that our porting effort evaluation mechanism captures this phenomenon.

It is also interesting to note that our porting effort values are not bounded by any upper limit. The porting effort could be arbitrarily large depending on the number and the complexity of the semantic predicates. We believe that this property of the metric reflects current practice. That is, it is always possible to find or create two new APIs E and F , such that the effort it takes to port an application from E to F is greater than any previously known upper bound. Our porting effort evaluation mechanism captures this property.

3.2 Publicly Available, Real-World APIs

We next investigate the efficacy of our approach using popular, publicly available web APIs. We list these APIs below. To evaluate our porting effort metric, we have augmented the APIs with semantic specifications manually. To enable this, we carefully analyze the API documentation and examples related to each of these web APIs. Specifically, we identify an important operation from each API set that was present across the set and specify its pre/postconditions using our specification language. Thus, these results pertain the similarity between an individual API operation that is common to all APIs in a set (either social media, airline services, or digital media).

- Social media login APIs: Facebook, Google, LlinkedIn, Twitter, Yahoo, Hi5, Amazon
- Airline itinerary search APIs: American Airlines, British Airways, Cathay Pacific, Delta Airlines, Emirates, Eti-

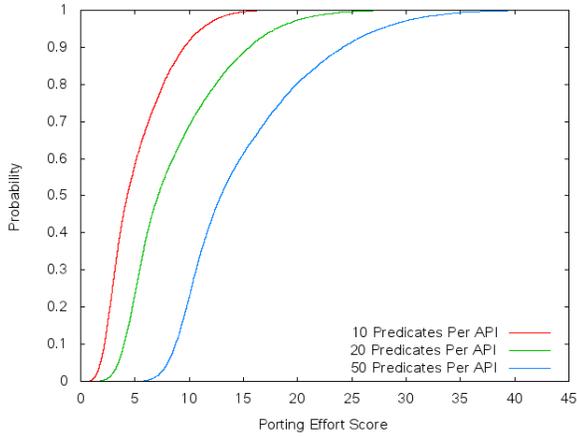


Figure 1: Porting effort CDFs for randomly generated APIs

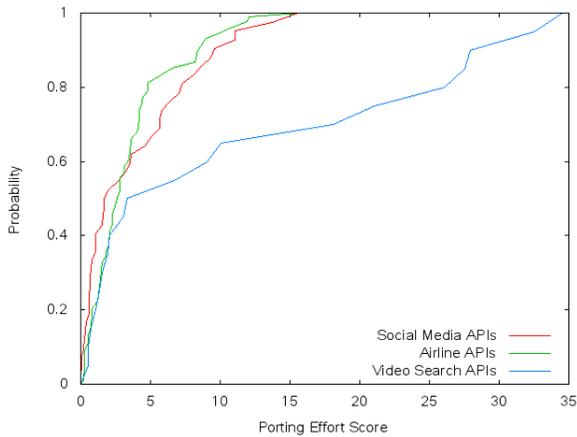


Figure 2: Porting Effort CDFs for real world APIs

had, Singapore Airlines, United Airlines, Virgin America

- Digital media video search APIs: Youtube, iTunes, MovieDB, RottenTomatoes, Vimeo

We then compute the porting effort among each pair of APIs within each of the above three categories. We present the CDFs of the results in figure 2.

The data shows that a fairly large proportion of the API pairs have a low porting effort. For instance, in all three populations (social media, airlines and video search), 50% of the pairs have a porting effort of 3.3 or less, a characteristic not present in the data obtained from the randomly generated APIs. This is because, unlike in the randomly generated populations where most APIs are completely unrelated to each other, in real world API populations most APIs can and do have commonalities. For instance, most social media login APIs have similar constraints on username and password. Most airline APIs have similar requirements with respect to specifying departure and arrival cities, travel dates and the number of passengers. Most video search APIs also exhibits similar constraints, in that most APIs at least

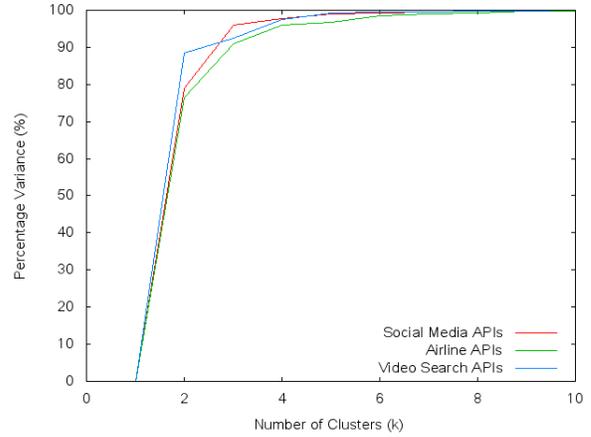


Figure 3: Percentage variance of porting effort

accept simple text queries to perform keyword-based search. These similarities simplify and make application porting between a number of API alternatives practical.

The CDFs of the social media APIs and the airline APIs follow relatively similar trends. However, the CDF of the video search APIs deviates from the other two and reaches a maximum porting effort value close to 35. A closer look at the API specifications showed that social media APIs and the airline APIs are similar in terms of their average semantic predicate count (8.1 and 9.3 respectively). For the video search APIs, the average predicate count is as high as 15.6 thus resulting in an increased porting effort among them. Also, some of the video search APIs have a large number of semantic predicates compared to the others. For instance, Youtube search API has 28 semantic predicates, and the iTunes search API has 30 semantic predicates. Therefore ports that involve these APIs tend to be much more complicated than the others.

3.3 Categorizing API Porting Difficulty

Given the observed properties of the methodology (particularly for the real-world APIs), one way which we believe it to be useful is in determining categories of difficulty. That is, it should be possible for the methodology to “cluster” API ports into groups that can be ranked in terms of difficulty (*e.g.* is a port “easy” or “hard”?)

To investigate this hypothesis we use k -means clustering to classify the results into two groups (*i.e.* $k = 2$). Figure 3 shows, for each sample set, the ratio of the variance explained by the categorization to the total variance in the set. Typically, this analysis shows an “elbow” in the curve corresponding to the point where further categorization adds little explanatory power. In our study, that point of diminishing returns appears at $k = 2$.

Thus, for these API operations, it appears that our methodology should be able to divide pairwise porting effort into two categories: “easy” and “difficult.” We then asked two of our lab members (lets call them D_1 and D_2) conversant with web services but not otherwise associated with this project to categorize the porting difficulty of a subset of the porting

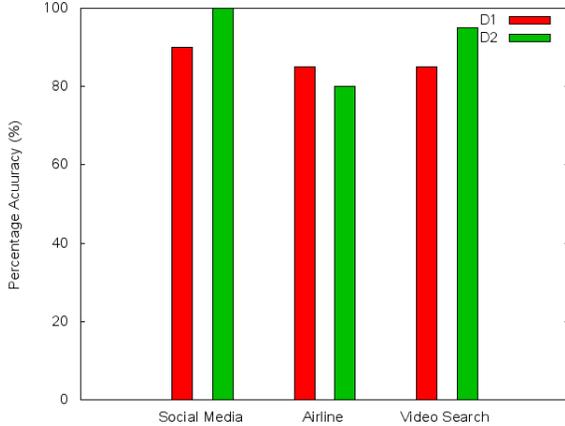


Figure 4: Percentage Accuracy of the Classification

possibilities in each set as either “easy” or “difficult.”

We gave these developers three sample sets, each consisting of 5 API specifications, randomly chosen from the above three categories (social media, airlines and video search). We then asked each developer to analyze the API specifications pairwise, and classify all possible pairs into two groups – easy and difficult – depending on the potential complexity of porting an application from one API to another. We also computed the porting effort between these web APIs using our own prototype, and used k -means clustering to classify the results into two groups (i.e. $k = 2$).

Figure 4 shows the percentage accuracy of the classifications computed using our formal mechanism with respect to the classifications provided by developers D_1 and D_2 respectively. We compute the percentage accuracy as the ratio of the number of entries classified as the same (i.e. agreement between the developer and the methodology) to the total sample size. In terms of a simple categorization, the agreement is good. Indeed, developer D_2 and the methodology obtained the same classification (100% accuracy) for the social media API operation.

3.4 Overhead

Finally, we evaluate the time overhead associated with computing web API porting effort using our mechanism. We employ our randomly generated set of 100 API specifications and compute the porting effort between each pair of APIs. We measure the time elapsed for all computations and then compute the average time per API pair. We repeat the experiment, varying the total number of semantic predicates in each API specification. We report the average times that we observe in these experiments in figure 5.

For web APIs with 10 semantic predicates, our evaluation method takes less than 10 ms. This increases up to 200 ms when the predicate count is increased to 50. This increase in execution time is due to the pairwise AST comparison operations performed by our algorithm. That is, when computing the porting effort between two APIs, our prototype compares each precondition of the source API against each precondition of the target API. In the same fashion, our

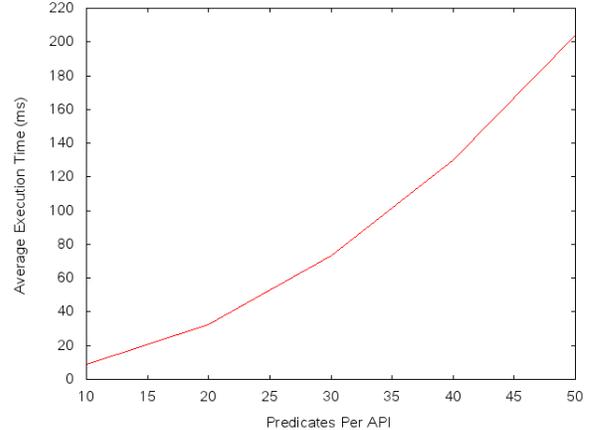


Figure 5: Average execution time of the porting effort evaluation algorithm

prototype compares each postcondition of the source API against each postcondition of the target API. Therefore the number of AST comparisons performed is polynomial in the number of semantic predicates. Hence the average execution time of our algorithm increases polynomially with the increase in semantic predicates. However, for web APIs with 10 semantic predicates, the average execution time is below 10 ms and for web APIs with 20 semantic predicates, the average execution time is well below 50 ms. Since most of the real world web APIs that we have studied to date have a small number of predicates (the max was 30), our approach is not likely to impose a significant time overhead on the development process for applications. If required, the algorithm can be easily parallelized by running the pairwise AST comparisons in parallel to reduce the overhead further.

Overall, our porting effort evaluation method produces useful results with a high level of accuracy. The method is efficient, and can be easily applied to real world web APIs. The Python-based syntax simplifies documentation and publication of API semantics (relative to semantic ontologies, state machines, and formal logic) by API providers. If an API provider fails to publish API semantics in our language, API consumers (developers) can easily create API specifications on their own by converting the semantics of API operations described in the API documentation into Python code.

4. RELATED WORK

Our work builds upon and extends research from a number of other areas in computer science. These areas include programming language and web service semantics, analysis and verification.

Our approach of using axiomatic semantics to describe web APIs is rooted in the work of Floyd [34] and Hoare [24]. Floyd modeled computer programs as digraphs where vertices represent program statements and edges represent control flow. Predicates representing correctness conditions are attached to the edges. Hoare introduced the notion of Hoare triples and constructed a formalism for reasoning about program correctness using them. A Hoare triple is a logical construct of the form $P\{C\}Q$ where C is a command (an

operation) in a program, P is the set of preconditions of C and Q is the set of postconditions of C . We adapted this formalism into our work where we reason about web services by describing their operations along with the respective preconditions and postconditions. Hoare’s seminal work on using axiomatic semantics to reason about program correctness excludes side effects and arbitrary procedure calls. In this work, we follow the same approach for semantic predicate description language to facilitate low complexity and thus fast evaluation and comparison of API porting effort.

Several researchers have been successful in using axiomatic semantics to reason about the correctness and behavior of software constructs. Hoare himself, along with Wirth showed how axiomatic semantics can be used to describe Pascal programs [35]. Fikes and McGuiness used axiomatic semantics to describe RDF data models [36]. Gegg-Harrison et al introduced ProVIDE [37], a software development tool that allows the user to establish program correctness via specifying postconditions and then generating the corresponding preconditions. Black used axiomatic semantics to verify the behavior of a secure web server [38].

Our guidelines for comparing web API semantics are loosely based on Hoare’s rule of consequence [24]. The rule of consequence states that if $P\{C\}Q$ and $P'\{C'\}Q'$ are two Hoare triples such that $P \Rightarrow P'$ and $Q' \Rightarrow Q$, then the command C' can be used in any context where the command C can be used. This is because C' has more permissive preconditions and more restrictive postconditions compared to C . We follow a similar rule when comparing web APIs with unequal number of preconditions or postconditions. Similar arguments have been discussed by Naumann [39] and Olderog [40] in the past.

The use of programming language syntax for expressing program semantics and contracts is a widely used concept. JML [26] uses two primary annotations (`requires` and `ensures`) to document the preconditions and postconditions of Java methods using Java syntax. Spec# [27] provides similar functionality for the C# language. SPARK language [41] has built-in contract documentation features, where contracts are encoded in the source code as Ada comments. These technologies use the documented semantics or contracts mostly for verification purposes. That is, they verify whether the program adheres to the given contract at the runtime. We use the documented semantics at the development time to reason about web service semantics and porting effort by applying static analysis methods.

The use of AST representations to compare programs and reason about them is also well researched. Our approach is heavily based on the work of Baxter et al, where they used AST comparison methods for detecting program clones [32]. Baxter et al introduced the notion of syntactic similarity (based on the Dice coefficient), as opposed to exact matches, as a more practical means of finding program segments with similar functionality and behavior. Cui et al showed how to use AST comparison methods for source code plagiarism detection [42]. They showed that AST comparison based methods are capable of finding a wide range of similarities between different programs. Hashimoto and Mori augmented AST comparison methods with heuristics-

driven techniques so that they can be used to efficiently analyze the differences between programs written in a wide range of programming languages [43]. Neamtiu et al used AST comparison methods to track down and analyze how a program code base has evolved over time [44].

Bianchini et al introduced the notion of semantics-enabled web API selection patterns [45]. One of the selection patterns they discuss is the substitution pattern which aims at finding a web API that can be used to substitute another API (i.e. porting). They presented a formalism to model and quantify this selection pattern based on semantic ontologies. However, constructing comprehensive semantic ontologies requires a lot of time and manual effort, and therefore such techniques are difficult to apply in practice. Also they mainly focus on semantically annotating the input/output data elements of web APIs, whereas we look at the functional and behavioral traits of the web APIs.

There is a large body of work that uses techniques like process models (e.g. [46]), state machine models (e.g. [30]) and logic (e.g. [11]) to reason about web service behavior. However these formalisms are aimed at addressing the issues such as discovery, monitoring and verification. Our work deviates from these formal methods, in the sense we attempt to reason about developer experience of different web services, in the sense how much effort a developer has to put in to port an application from one web API to another.

5. CONCLUSIONS

Increasingly, web, mobile, and cloud developers integrate publicly available web services, exposed via well-defined APIs, into their Internet accessible applications. Doing so simplifies and expedites software development, testing, deployment, and management of these applications. Despite these benefits that arise from decoupling of APIs from the implementations they serve, the web service model has also introduced a key challenge for developers: API churn – constant API evolution (versioning) and emergence of alternative, competitive implementations for the same API. As a result, it is critical that developers be able to efficiently estimate and reason about the work required to migrate their applications from one API (or API version) to another.

In this paper, we investigate a new methodology for automatically quantifying such *porting effort*. Our approach defines a simple language based on Python with which API developers specify the semantics of API operations, a tool set that consumes and analyzes specified API semantics, and a new metric, based on a method used to detect code plagiarism, to estimate porting effort between pairs of API operations. We evaluate a prototype of this approach using randomly generated APIs to measure the sensitivity to the parameters we employ, and using competitive, publicly available APIs to determine its efficacy on real-world APIs. We employ k -means clustering to divide API ports into groups that can be ranked in terms of difficulty and show that the variance is explained by 2-3 clusters. Finally, we show that computation of our porting effort metric introduces minimal overhead, making it sufficiently practical to include in a developer’s tool chain.

6. REFERENCES

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